AProVE (KoAT + LoAT)

Automatic Termination Analysis of C Programs

Nils Lommen, Florian Frohn, and Jürgen Giesl

Overview

- AProVE (KoAT + LoAT) [1] is a framework to analyze termination of C Programs
- Programs are transformed into Integer Transition Systems (ITSs)
- ITSs are analyzed by our tools KoAT [2] and LoAT [3]



LLVM Program

- C program is compiled into LLVM code using Clang.
- LLVM fragment of the loop body:

1	%10 = load %1	# load ×
2	%11 = mul 3 %10	∦ multiply x by 3
3	store %11, %1	# store x
4	$\%12 = load \ \%2$	# load y
5	$\%13 = mul \ 2 \ \%13$	∦ multiply y by 2
6	store %13, %2	# store y
7	br %6	# jump to loop guard

Symbolic Execution Graph (SEG) & ITS

SEG represents all possible program runs, augmented with invariants:

- Its nodes are *abstract states* that represent sets of actual program states
- SEG handles the heap, pointer arithmetic, and recursive data structures
- $\bullet~\mbox{LLVM}$ code is transformed automatically into an SEG

ITSs are a simple language for integer programs:

- Turing-complete formalism with only integer variables over $\ensuremath{\mathbb{Z}}$
- SEG is transformed into ITS

$\rightarrow \underbrace{\ell_0}_{t_0} \underbrace{t_0}_{\ell_1} \underbrace{\ell_1}_{\eta(x)} \underbrace{t_1 : \varphi = (x < y)}_{\eta(x)} \\ \eta(x) = 3 \cdot x \\ \eta(y) = 2 \cdot y$

- KoAT (Termination & Upper Time Bounds)
- Automated complexity and termination analysis of ITSs
- Alternating modular inference of runtime and size bounds
- How often can a transition be executed?
 - Multiphase Linear Ranking Functions
 → Use SMT-solver Z3 to infer well-founded relation
 - TWN-Loops
 → Reduce termination problem to SMT problem
 - Completeness for the class of so-called TWN-loops
- How *large* are the variables?
 - Compute bounds for each change of a variable
 → Over-approximate the number of changes by runtime bounds
 - Use runtime bounds and closed forms of loops

LoAT (Non-Termination and more)

Features

- Techniques
- non-termination ADCL DFS + acceleration
- *lower* time bounds **ABMC** BFS + acceleration
- safety / unsafety **TRL** BFS + recurrence analysis

Non-term. via <u>A</u>cceleration <u>D</u>riven <u>C</u>lause <u>L</u>earning

- Depth-first exploration of state space
- Applies <u>acceleration</u> when a loop is encountered under-approximation of the loop's transitive closure
- Non-term. proofs as "by-product" of acceleration
- Exploits *redundancy* to cut off infinite branches



References

[2]

- [1] Nils Lommen and Jürgen Giesl. AProVE (KoAT + LoAT) Website: https://koat.verify.rwth-aachen.de/svcomp25.
 - Nils Lommen, Éléanore Meyer, and Jürgen Giesl. KoAT Website: https://koat.verify.rwth-aachen.de/.
- [3] Florian Frohn and Jürgen Giesl. LoAT Website: https://loat-developers.github.io/LoAT/.





Exemplary C Program

int main() {

Does the following program terminate?

extern int _nondet(void);

int x = _nondet();

int y = _nondet();

while (x < y) { x = 3*x;

y = 2*y;

return 0;